

# Code Converters

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# Binary To Gray Code Converter

Gray Code system is a binary number system in which every successive pair of numbers differs in only one bit. It is used in applications in which the normal sequence of binary numbers generated by the hardware may produce an error or ambiguity during the transition from one number to the next.

For example, the states of a system may change from 3(011) to 4(100) as- 011 — 001 — 101 — 100. Therefore there is a high chance of a wrong state being read while the system changes from the initial state to the final state.

This could have serious consequences for the machine using the information. The Gray code eliminates this problem since only one bit changes its value during any transition between two numbers.

# Binary to Gray Code converter Truth table

Binary				Gray Code			
b <sub>3</sub>	b <sub>2</sub>	b <sub>1</sub>	b <sub>0</sub>	g <sub>3</sub>	g <sub>2</sub>	g <sub>1</sub>	g <sub>0</sub>
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	0	0	0	1	1
0	0	1	1	0	0	1	0
0	1	0	0	0	1	1	0
0	1	0	1	0	1	1	1
0	1	1	0	0	1	0	1
0	1	1	1	0	1	0	0
1	0	0	0	1	1	0	0
1	0	0	1	1	1	0	1
1	0	1	0	1	1	1	1
1	0	1	1	1	1	1	0
1	1	0	0	1	0	1	0
1	1	0	1	1	0	1	1
1	1	1	0	1	0	0	1
1	1	1	1	1	0	0	0

Let  $b_3 b_2 b_1 b_0$  be the bits representing the binary numbers, where  $b_0$  is the LSB and  $b_3$  is the MSB, and  
Let  $g_3 g_2 g_1 g_0$  be the bits representing the gray code of the binary numbers, where  $g_0$  is the LSB and  $g_3$  is the MSB.

To find the corresponding digital circuit, we will use the K-Map technique for each of the gray code bits as output with all of the binary bits as input.

# K map for g0

		b1,b0			
		00	01	11	10
b3,b2	00	0	1	0	1
	01	0	1	0	1
	11	0	1	0	1
	10	0	1	0	1

$$g_0 = b_0 b'_1 + b_1 b'_0 = b_0 \oplus b_1$$

# K map for g1

		b1,b0			
		00	01	11	10
b3,b2	00	0	0	1	1
	01	1	1	0	0
	11	1	1	0	0
	10	0	0	1	1

$$g_1 = b_2 b_1' + b_1 b_2' = b_1 \oplus b_2$$

# K map for g2

		b1,b0			
		00	01	11	10
b3,b2	00	0	0	0	0
	01	1	1	1	1
	11	0	0	0	0
	10	1	1	1	1

$$g_2 = b_2b'_3 + b_3b'_2 = b_2 \oplus b_3$$

# K map for g3

		b1,b0			
		00	01	11	10
b3,b2	00	0	0	0	0
	01	0	0	0	0
	11	1	1	1	1
	10	1	1	1	1

$$g_3 = b_3$$

# Corresponding minimized boolean expressions for gray code bits

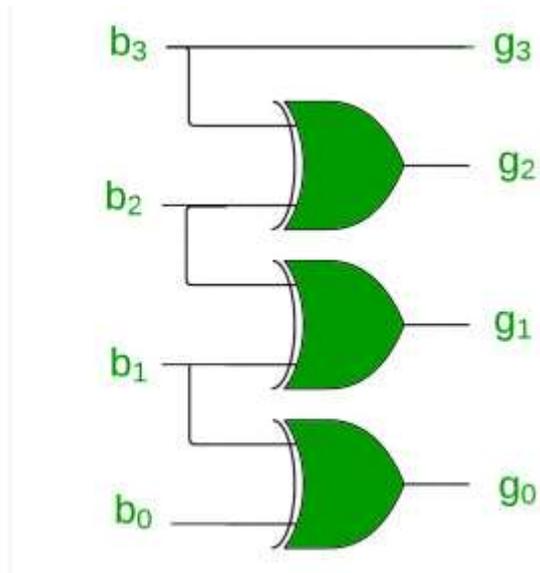
$$g_0 = b_0b'_1 + b_1b'_0 = b_0 \oplus b_1$$

$$g_1 = b_2b'_1 + b_1b'_2 = b_1 \oplus b_2$$

$$g_2 = b_2b'_3 + b_3b'_2 = b_2 \oplus b_3$$

$$g_3 = b_3$$

# Code converter



# BCD to excess-3 code converter

As is clear by the name, a BCD digit can be converted to its corresponding Excess-3 code by simply adding 3 to it.

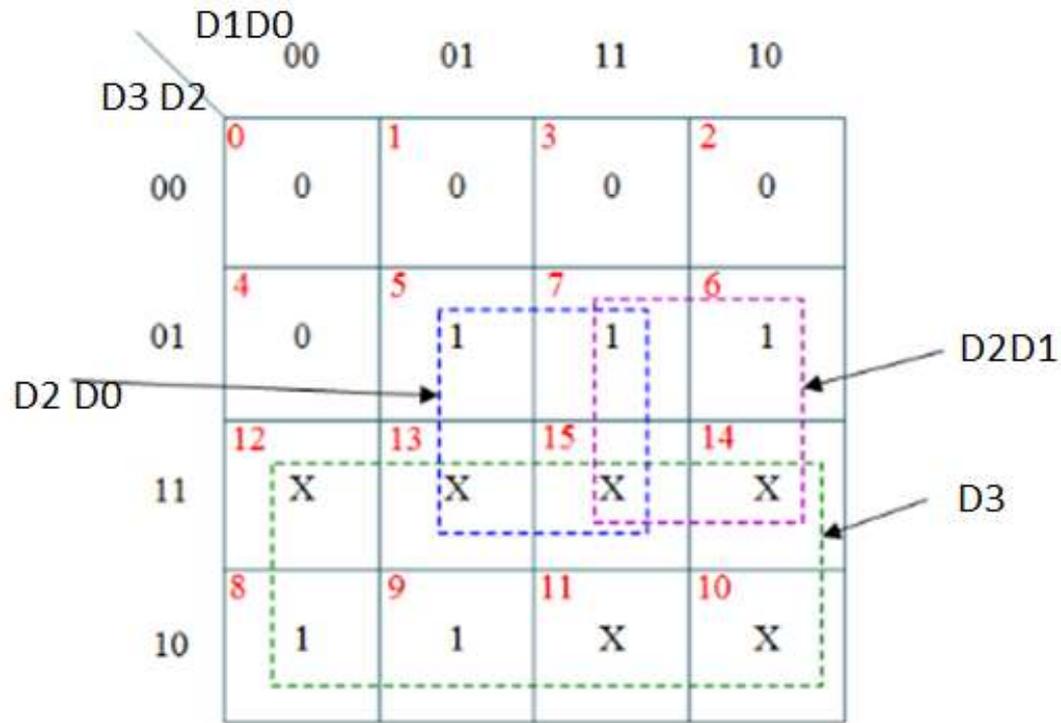
Let  $A, B, C,$  and  $D$  be the bits representing the binary numbers, where  $D$  is the LSB and  $A$  is the MSB, and

Let  $w, x, y,$  and  $z$  be the bits representing the gray code of the binary numbers, where  $z$  is the LSB and  $w$  is the MSB.

The truth table for the conversion is given below. The X's mark don't care conditions.

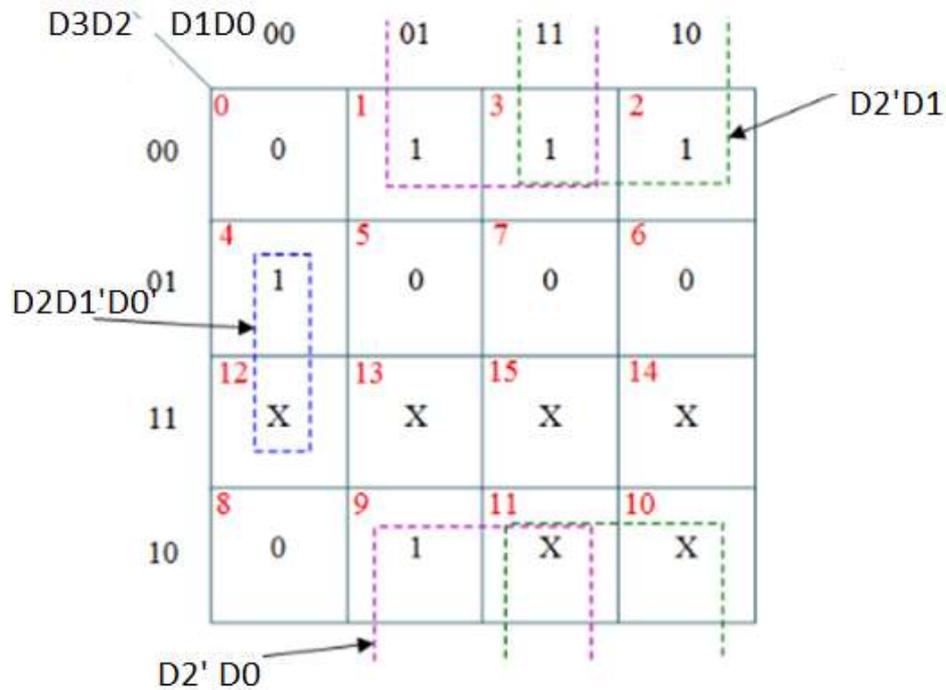
BCD(8421)				Excess-3			
D3	$\bar{D}2$	$\bar{D}1$	D0	E3	E2	$\bar{E}1$	E0
0	0	0	0	0	0	1	1
0	0	0	1	0	1	0	0
0	0	1	0	0	1	0	1
0	0	1	1	0	1	1	0
0	1	0	0	0	1	1	1
0	1	0	1	1	0	0	0
0	1	1	0	1	0	0	1
0	1	1	1	1	0	1	0
1	0	0	0	1	0	1	1
1	0	0	1	1	1	0	0
1	0	1	0	X	X	X	X
1	0	1	1	X	X	X	X
1	1	0	0	X	X	X	X
1	1	0	1	X	X	X	X
1	1	1	0	X	X	X	X
1	1	1	1	X	X	X	X

# K map for E3



$$E3 = D3 + D2D0 + D2D1$$

# K map for E2



$$E2 = D2'D1 + D2'D0 + D2D1'D0'$$

# K map for E1

D3 D2		D1 D0			
		00	01	11	10
00	0	1	3	2	
01	4	5	7	6	
11	12	13	15	14	
10	8	9	11	10	

The Karnaugh map shows the following values in the cells:

D3 D2 \ D1 D0	00	01	11	10
00	1	0	1	0
01	1	0	1	0
11	X	X	X	X
10	1	0	X	X

Annotations:

- A blue dashed box highlights the 1s in the first column (D1=0, D0=0), labeled  $D1'D0'$ .
- A green dashed box highlights the 1s in the third column (D1=1, D0=0), labeled  $D1D0$ .

$$E1 = D1D0 + D1'D0'$$

# K map for z

D1D0 \ D3D2	00	01	11	10
00	0 1	1 0	3 0	2 1
01	4 1	5 0	7 0	6 1
11	12 X	13 X	15 X	14 X
10	8 1	9 0	11 X	10 X

$E0 = D0'$

$D0'$

# Minimized expressions

$$E3 = D3 + D2D0 + D2D1$$

$$E2 = D2'D1 + D2'D0 \\ + D2D1'D0'$$

$$E1 = D1D0 + D1'D0'$$

$$E0 = D0'$$

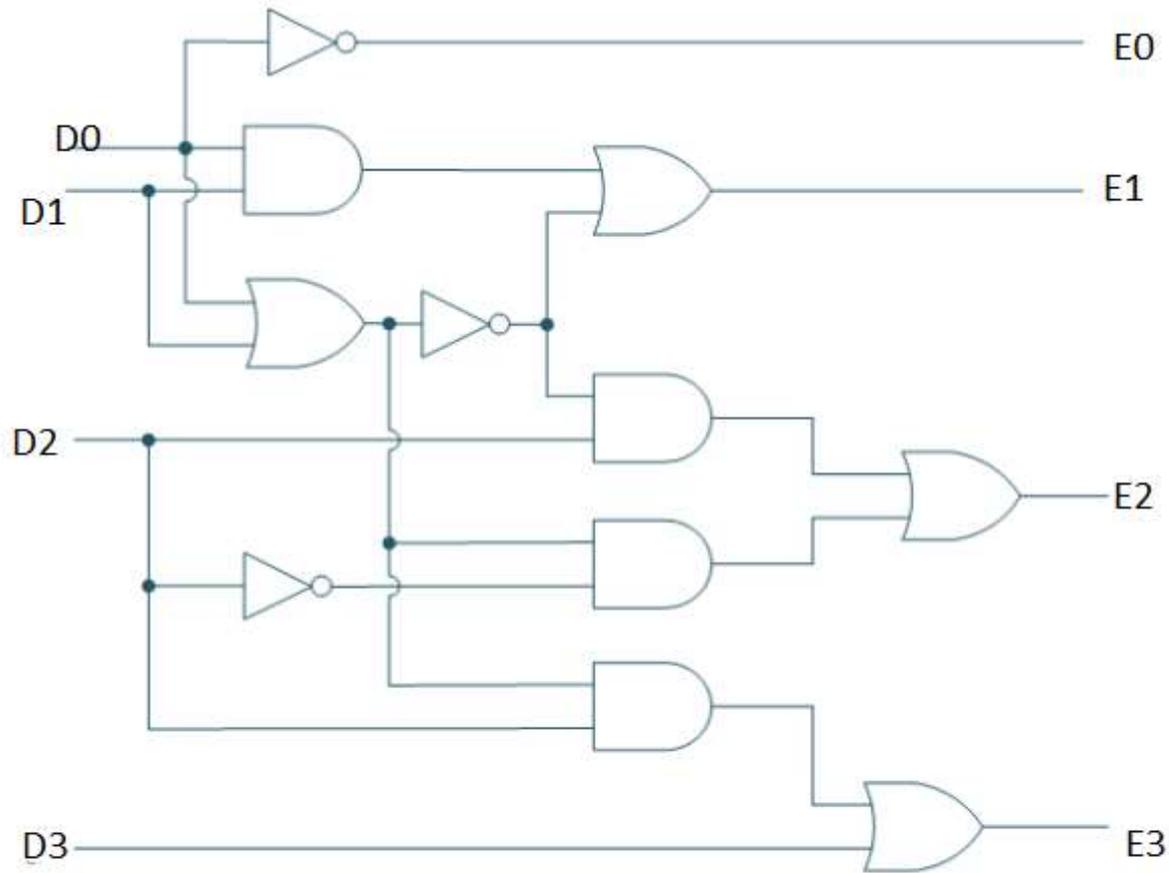
$$E3 = D3 + D2(D0 + D1)$$

$$E2 = D2'(D1 + D0) + D2D1'D0'$$

$$E1 = D0 \text{ EXNOR } D1$$

$$E0 = D0'$$

# BCD to Excess-3 code converter



# Practice Problem

## BCD to Gray code converter

# Truth table

BCD code				Gray code			
B <sub>3</sub>	B <sub>2</sub>	B <sub>1</sub>	B <sub>0</sub>	G <sub>3</sub>	G <sub>2</sub>	G <sub>1</sub>	G <sub>0</sub>
0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	1
0	0	1	0	0	0	1	1
0	0	1	1	0	0	1	0
0	1	0	0	0	1	1	0
0	1	0	1	0	1	1	1
0	1	1	0	0	1	0	1
0	1	1	1	0	1	0	0
1	0	0	0	1	1	0	0
1	0	0	1	1	1	0	1

# K map simplification

**For  $G_0$**

$B_3B_2 \backslash B_1B_0$	00	01	11	10
00	0	1	0	1
01	0	1	0	1
11	X	X	X	X
10	0	1	X	X

$$G_0 = \bar{B}_1B_0 + B_1\bar{B}_0$$

$$= B_1 \oplus B_0$$

**For  $G_1$**

$B_3B_2 \backslash B_1B_0$	00	01	11	10
00	0	0	1	1
01	1	1	0	0
11	X	X	X	X
10	0	0	X	X

$$G_1 = B_2\bar{B}_1 + \bar{B}_2B_1$$

$$= B_2 \oplus B_1$$

**For  $G_2$**

$B_3B_2 \backslash B_1B_0$	00	01	11	10
00	0	0	0	0
01	1	1	1	1
11	X	X	X	X
10	1	1	X	X

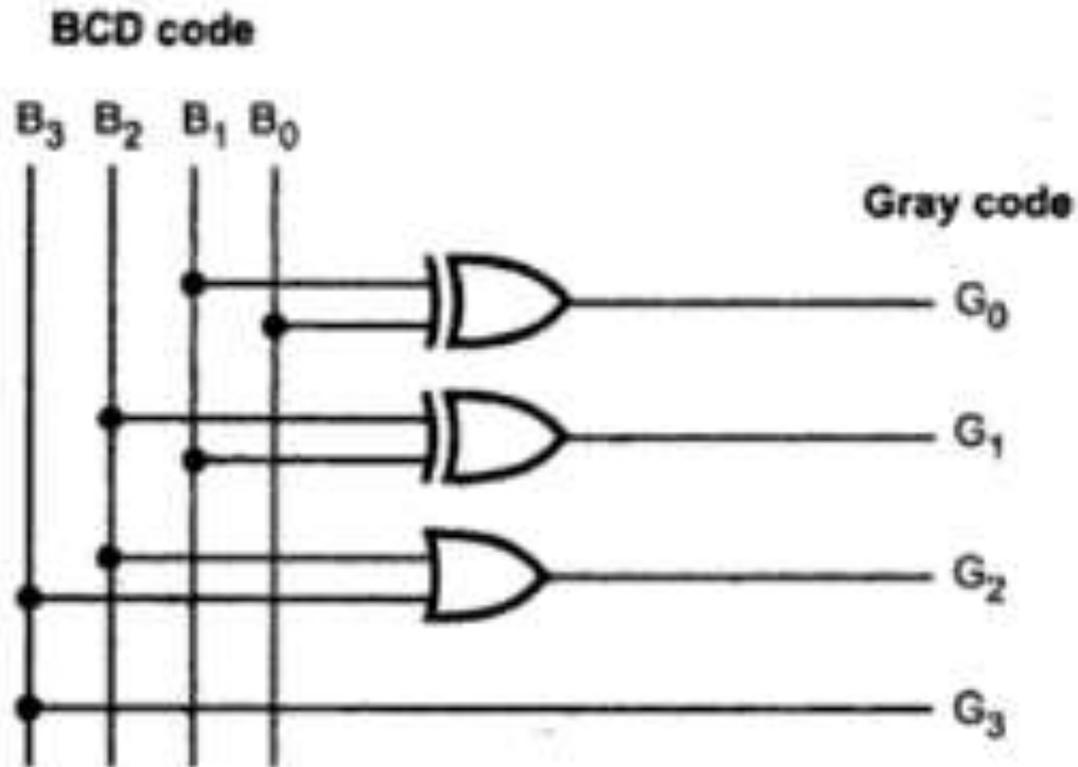
$$G_2 = B_2 + B_3$$

**For  $G_3$**

$B_3B_2 \backslash B_1B_0$	00	01	11	10
00	0	0	0	0
01	0	0	0	0
11	X	X	X	X
10	1	1	X	X

$$G_3 = B_3$$

# Circuit diagram



Thank you